

Linear contextual bandits with global constraints

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Based on joint work with Nikhil R. Devanur.

Application: Revenue management in internet advertising

- ▶ Operating delivery of ads so that long term revenue from the business is maximized
- ▶ Multi-billion dollar annual revenues

The screenshot shows a web browser displaying the MSN Lifestyle page. At the top, there is a green banner for Vaseline aloe fresh with the text "Vote on Your Favorite 'Lighten Up For Summer' Tip to Win a Trip to Malibu!" and a yellow button that says "For a chance to WIN" with a "CLICK HERE!" link. Below the banner is a search bar with "Search Lifestyle" and the Bing logo. The page content includes a "DECOR & ORGANIZING" section with a video player showing a woman in a blue and white patterned top and yellow gloves. To the right of the video is an advertisement for "ARM & HAMMER PLUS OXICLEAN POWER GEL" with a "CLICK FOR \$1.00 OFF" button. The page also features a navigation menu with categories like "Home", "Beauty & Fashion", "Decor & Organizing", "Live & Relationships", "Your Life", "Cooking", "Dating & Personal", "Go", "Health & Fitness", "Horoscopes", and "Travel".

Pay-per click advertising

Advertisers specify target user profiles, payment per click

- ▶ user opens a page at time t , matches target profile of many ads
- ▶ pick one ad
- ▶ “if the user clicks” on the shown ad, publisher gets paid

Uncertainty in future user profiles, uncertainty in clicks

“Click-through rate” depends on a combination of user profile and ad features.

Linear regression model

Click-through rates as a linear function of user and ad features.

- ▶ Let $x_{t,a}$ be a vector of features of (user t , ad a) combination
- ▶ On serving ad a to the user t , the chances of getting clicked is $w^T x_{t,a}$ for some unknown vector w .

Linear contextual bandit problem: explore-exploit in the feature space to learn w quickly.

Linear contextual bandits

In every round t , pick one of the many options (arms) in set A_t .

- ▶ For every $a \in A_t$, observe “context vector” $x_{t,a} \in \mathbb{R}^d$ before making the choice.
- ▶ On picking option a , observe reward $r_t \in [0, 1]$

Stochastic assumptions

- ▶ Reward r_t on picking arm a is i.i.d. from distribution with mean $w^T x_{t,a}$, w is unknown.
- ▶ No assumptions on the set A_t or context vectors – could be adversarial

Linear contextual bandits

Goal

- ▶ maximize sum of rewards $\sum_t r_t$
- ▶ minimize expected regret: compared to best context-dependent policy

$$\mathcal{R}(T) = \sum_t \max_{a \in A_t} w^T x_{t,a} - \mathbb{E}[\sum_t r_t]$$

UCB algorithms

- ▶ maintain a confidence ellipsoid around least-square estimate of w , use the most optimistic value \tilde{w}_t in the ellipsoid at time t
- ▶ at step t , play $\arg \max_{a \in A_t} \tilde{w}_t^T x_{t,a}$.
- ▶ achieve $\tilde{O}(d\sqrt{T})$ regret

Further considerations

Budget constraints!

Maximize the total value while not exceeding the budgets

$$\begin{aligned} &\text{maximize} && \sum_{t,a \in A_t} r_{t,a} y_{t,a} \\ & && \forall t, \quad \sum_{a \in A_t} y_{t,a} \leq 1 \\ & && \forall \text{ads } a, \quad \sum_{t: a \in A_t} r_{t,a} y_{t,a} \leq B_a \end{aligned}$$

Benchmark: Optimal context dependent policy?

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Now:

- ▶ Even if you know w , the choice at every step is not obvious
- ▶ Ad a or a' ?
 - ▶ Ad a has highest immediate revenue, but it appears in A_t very frequently
 - ▶ Ad a' has smaller immediate revenue, but there may not be another opportunity to use its budget.

Stochastic assumption and Benchmark

Stochastic assumption on A_t :

- ▶ Set A_t of context vectors is generated i.i.d. from some distribution \mathcal{D} over collection of sets of context vectors

Benchmark:

Value of best static context-dependent policy $q : A \rightarrow \Delta^N$,

$$\text{OPT} = \max_q \mathbb{E}[\sum_{t,a \in A_t} r_{t,a} q(A_t)_a] \\ \forall \text{ads } a, \mathbb{E}[\sum_{t:a \in A_t} r_{t,a} q(A_t)_a] \leq B_a$$

- ▶ Expectation over distribution of A_t s, and of $r_{t,a}$ given $w, x_{t,a}$.
- ▶ OPT is as good as any adaptive solution that knows w AND the distribution of A_t s.

Further considerations

- ▶ Multiple types of feedback – revenue, relevance, cost of serving, click, conversions, demographic targeting
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- ▶ Nonlinear
 - ▶ Risk on over-spend, under-delivery
 - ▶ Diversity of user profiles
 - ▶ Smooth delivery

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Can be modeled as convex constraints and objective

$$\begin{aligned} \max \quad & f(\sum_{t,a} \mathbf{v}_{t,a} y_{t,a}) \\ & \sum_{t,a} \mathbf{v}_{t,a} y_{t,a} \in \mathcal{S} \\ \forall t, \quad & \sum_a y_{t,a} \leq 1 \end{aligned}$$

Online decisions with unknown distribution of $\mathbf{v}_{t,a}$!

Linear contextual bandits with global convex constraints and objective

In every round t , pick one of the many options (arms) in set A_t .

- ▶ For every $a \in A_t$, observe “context vector” $x_{t,a} \in \mathbb{R}^d$ before making the choice.
- ▶ On pulling arm a , observe vector $\mathbf{v}_t \in [0, 1]^d$

Stochastic assumptions:

- ▶ Given that arm a is pulled, vector \mathbf{v}_t is i.i.d. from distribution with mean $W^T x_{t,a}$, matrix W is unknown.
- ▶ Set A_t of context vectors is generated i.i.d. from some distribution over collection of context vectors

Linear contextual bandits with global convex constraints and objective

Goal:

- ▶ Maximize $f(\frac{1}{T} \sum_{t=1}^T \mathbf{v}_t)$ while ensuring $\frac{1}{T} \sum_{t=1}^T \mathbf{v}_t \in \mathcal{S}$
- ▶ Minimize expected regret:

$$\text{Regret in Objective} = \text{OPT} - f\left(\frac{1}{T} \sum_{t=1}^T \mathbf{v}_t\right)$$

OPT is the value of best context-dependent policy (?)

$$\text{Regret in constraints} = d\left(\frac{1}{T} \sum_t \mathbf{v}_t, \mathcal{S}\right)$$

$d(\cdot, \cdot)$ is a distance function, e.g. L_1 distance.

Benchmark

Value of best static context-dependent policy

$$\text{OPT} = \max_q \left(f \left(\mathbb{E} \left[\left(\sum_{t,a} W^T x_{t,a} \right) q(A_t) \right] \right) \right) \text{ such that}$$
$$\mathbb{E} \left[\left(\sum_{t,a} W^T x_{t,a} \right) q(A_t) \right] \in S$$

- ▶ OPT is as good as any adaptive solution that knows W AND the distribution of contexts.

Our results

- ▶ $\tilde{O}(dT^{-1/3})$ regret bounds in both objective and distance from constraint set
- ▶ $\tilde{O}(d/\sqrt{T})$ regret bound if
 - ▶ value of OPT is known to sufficient accuracy.
 - ▶ concave objective, no constraints
 - ▶ only constraints: feasibility problem
- ▶ Important: no dependence on number of arms (possible user+ad types, which is exponential in d)

Main components of the algorithm

Handling unknown W

- ▶ On making an observation, update estimate of W using standard linear contextual bandit techniques

Handling uncertainty in contexts: Even with an accurate W , the problem is difficult: “online stochastic convex programming” [Agrawal, Devanur, SODA 2015].

Overview of the algorithm for known W

One dimensional problem, A_t of size 2, objective only.
(W.l.o.g. expected reward $w x_{t,a}$ can be replaced by $x_{t,a}$.)

At time t ,

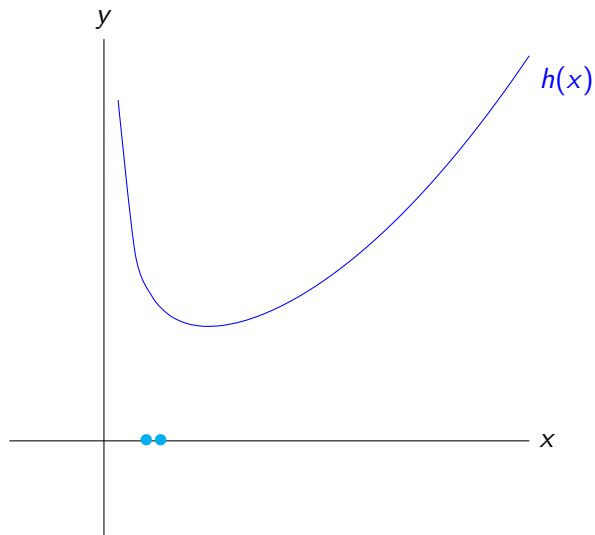
- ▶ you see random points $\{x_{t1}, x_{t2}\}$ on x -axis (stochastic assumption).
- ▶ Choose one of those points as x_t^\dagger .

Overall goal is to minimize $h(\frac{1}{T} \sum_{t=1}^T x_t^\dagger)$, where h is convex.

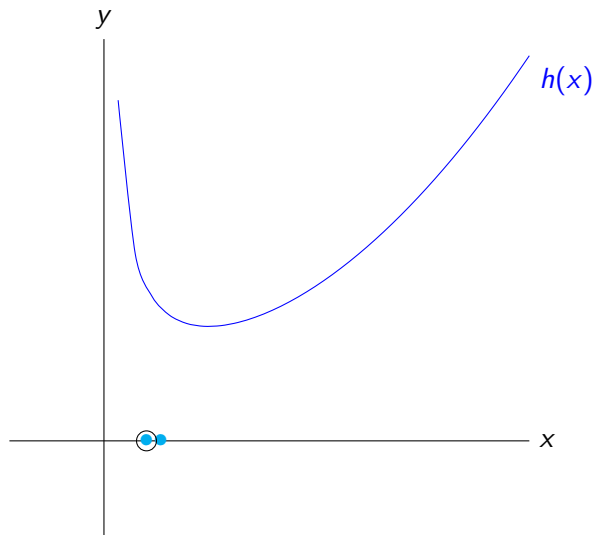
Regret

$$\mathcal{R}(T) = h\left(\frac{1}{T} \sum_{t=1}^T x_t^\dagger\right) - h\left(\frac{1}{T} \sum_{t=1}^T x_t^*\right).$$

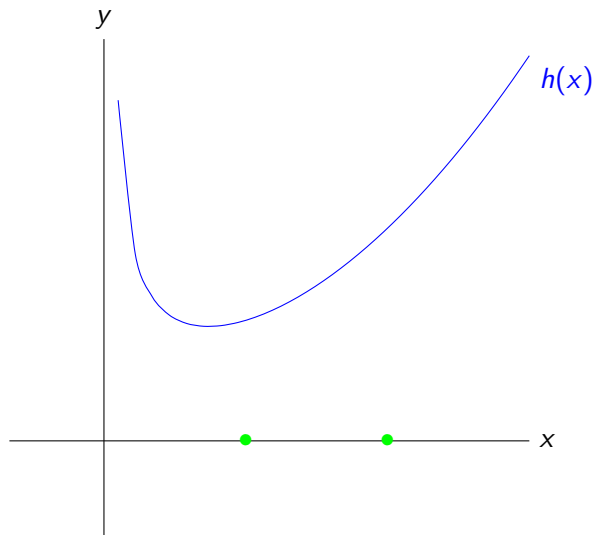
Overview by example



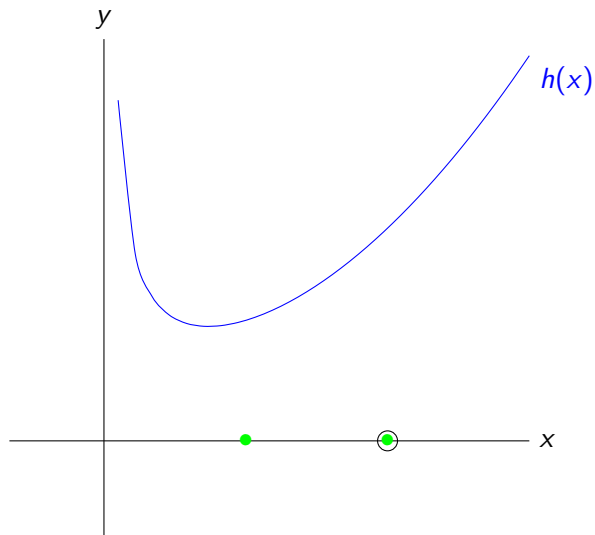
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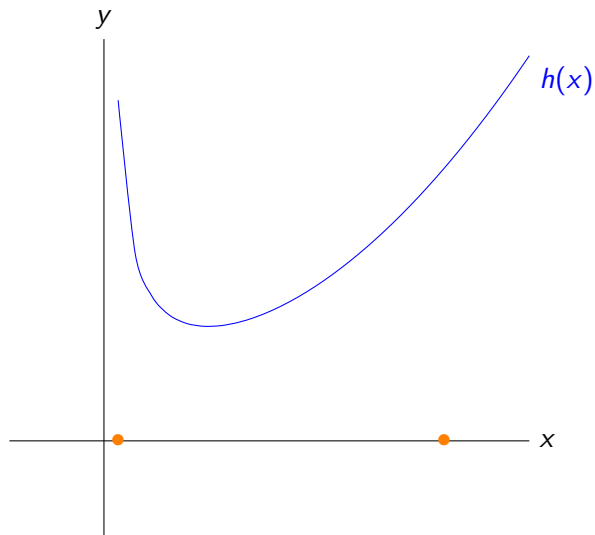
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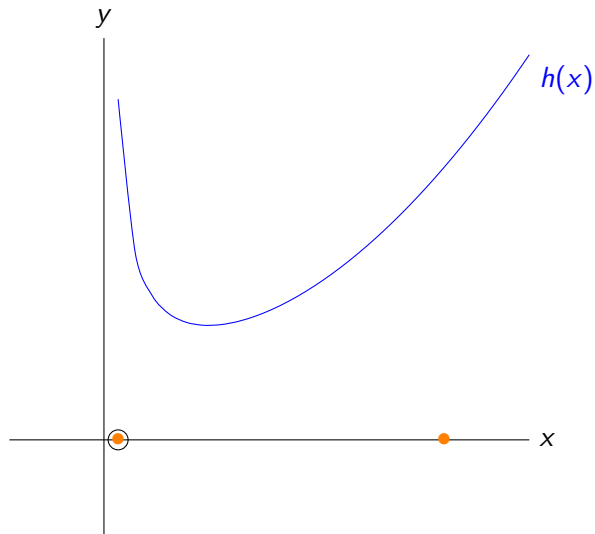
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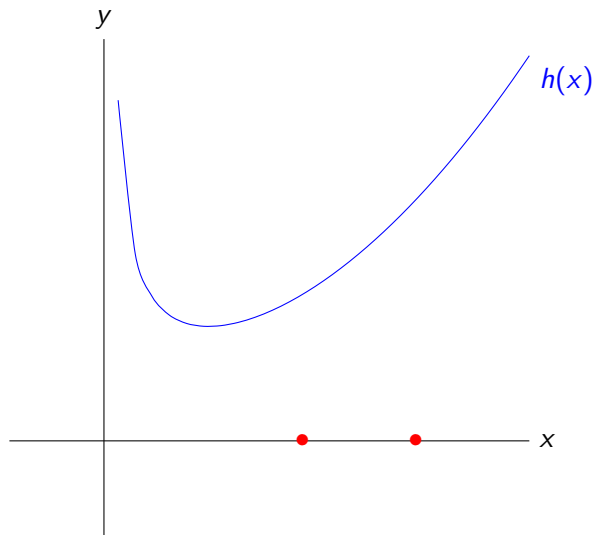
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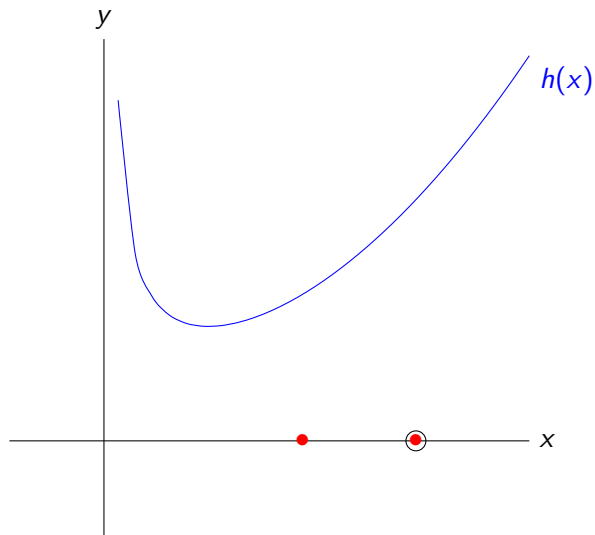
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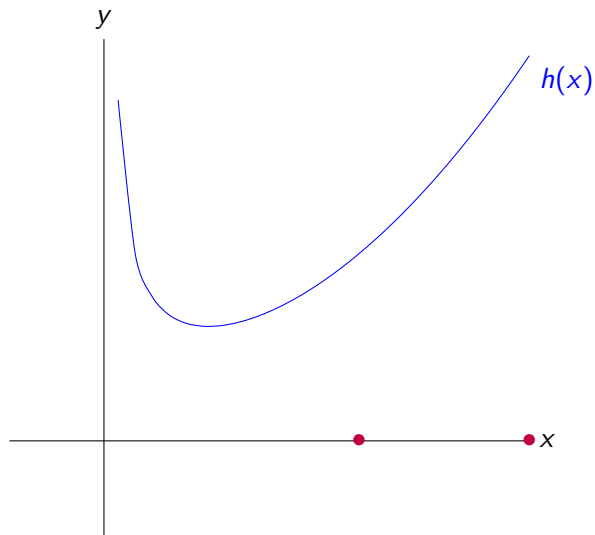
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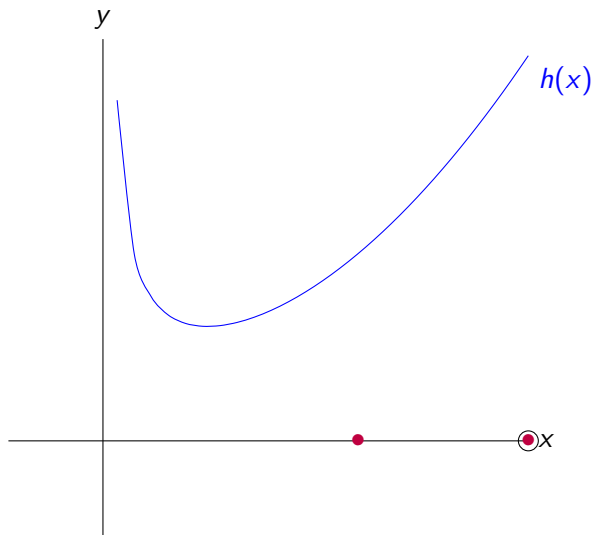
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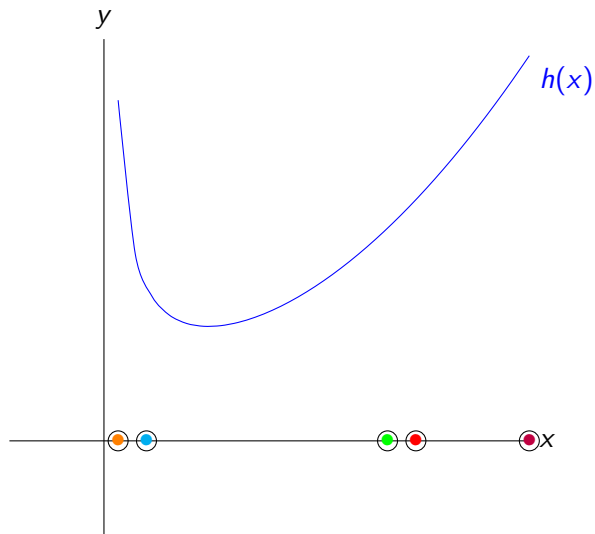
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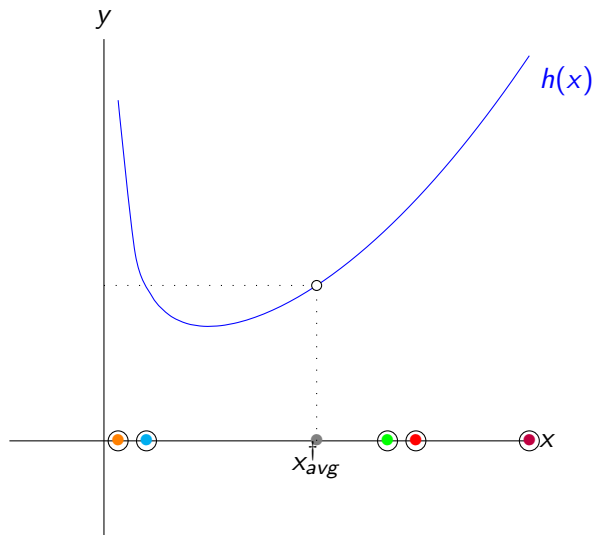
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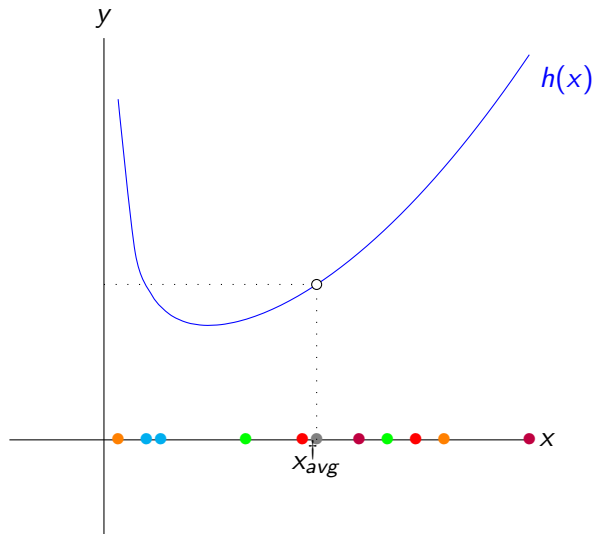
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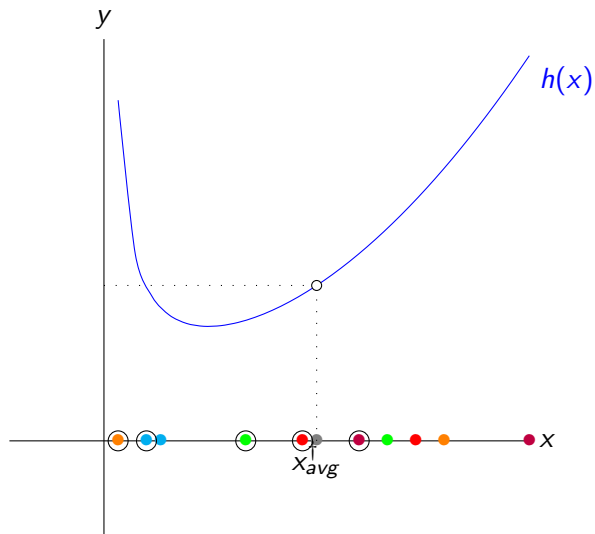
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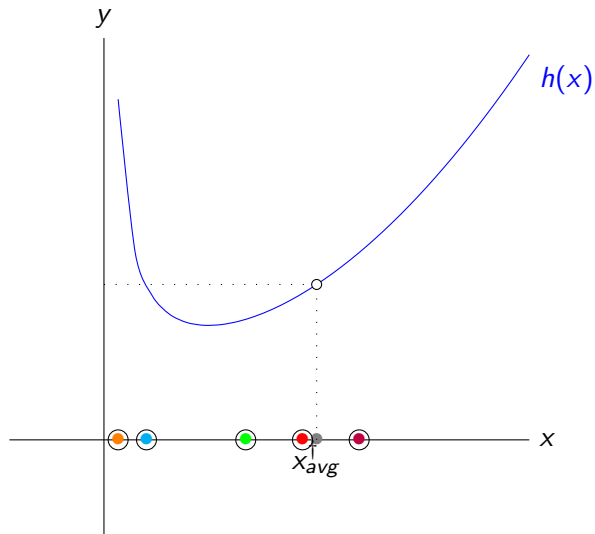
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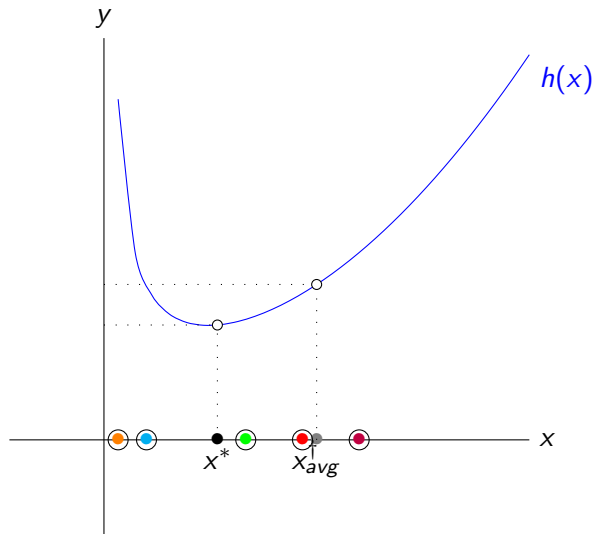
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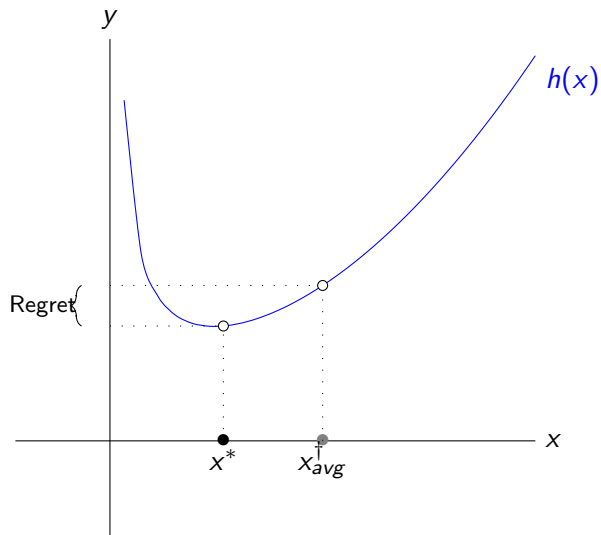
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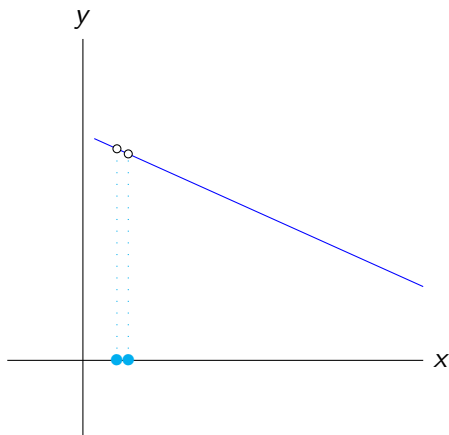
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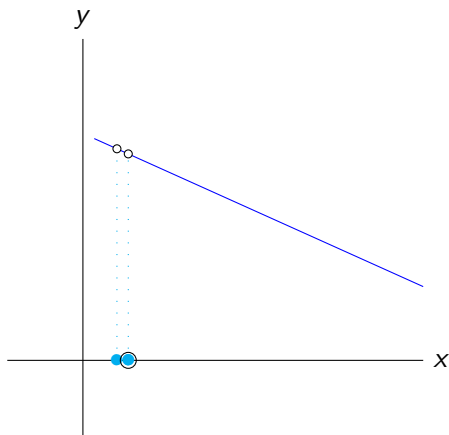


The simpler linear case



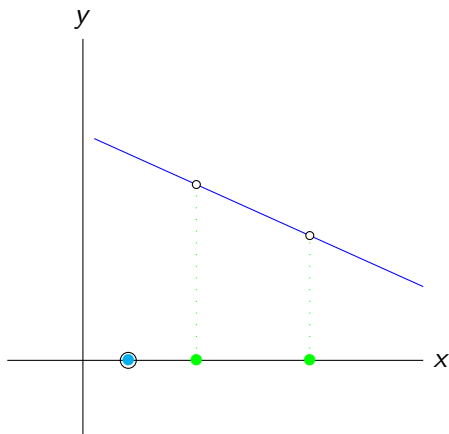
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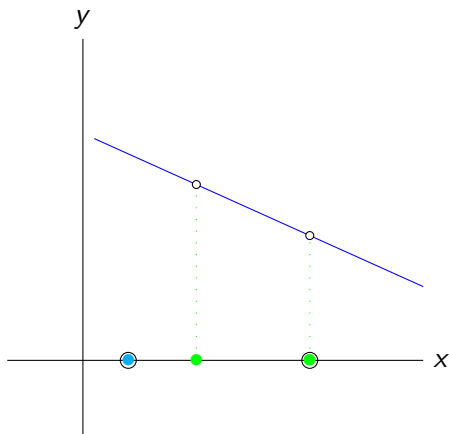
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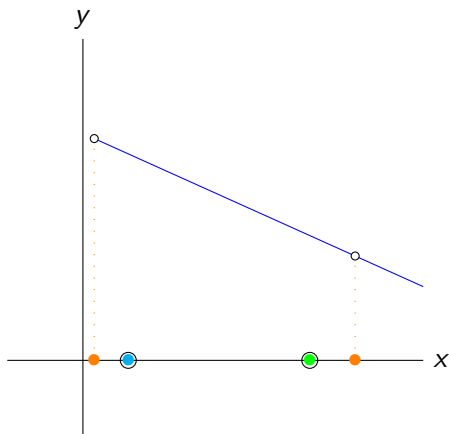
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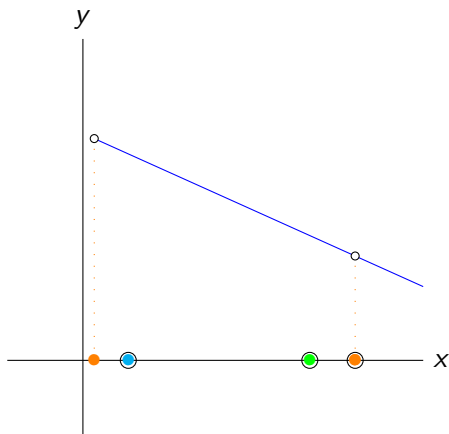
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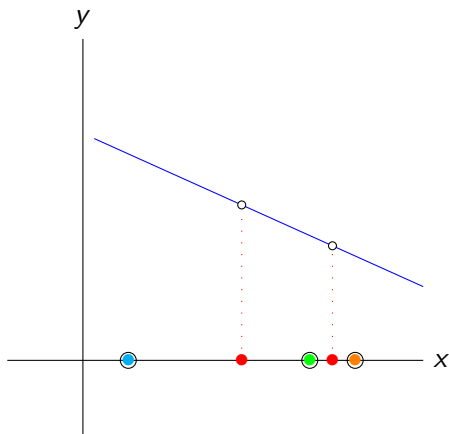
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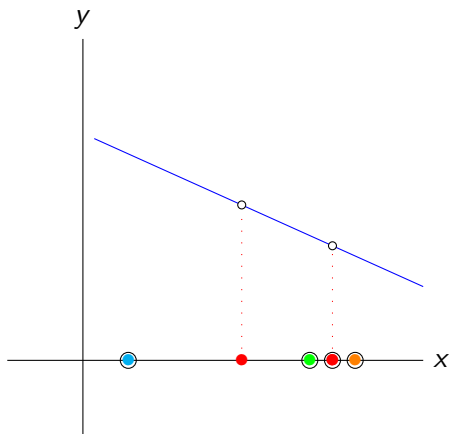
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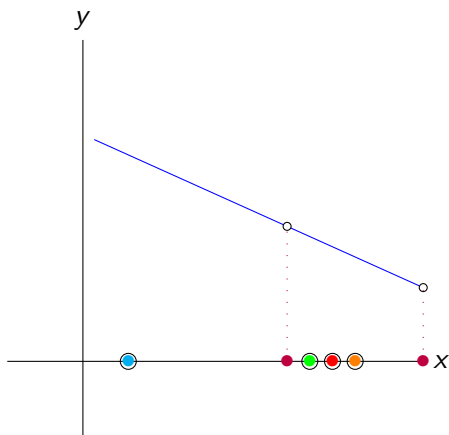
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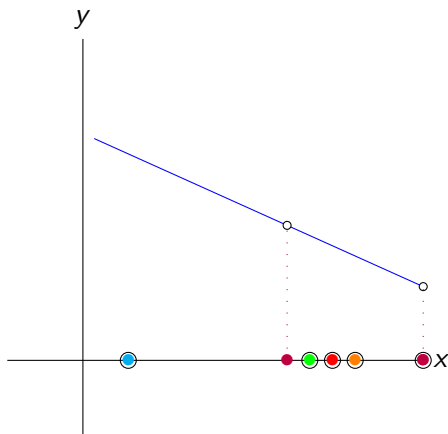
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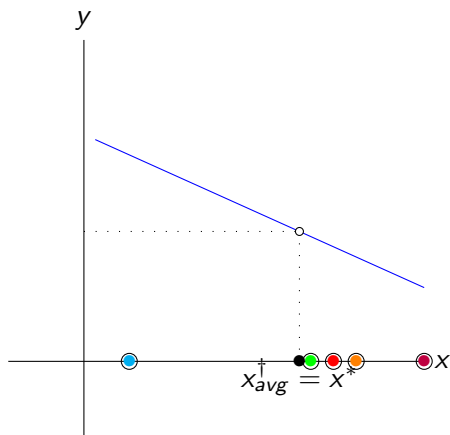
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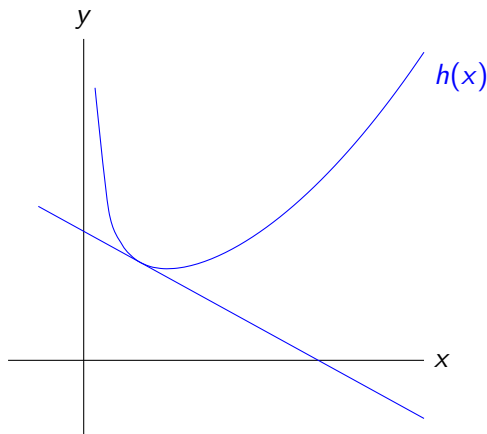
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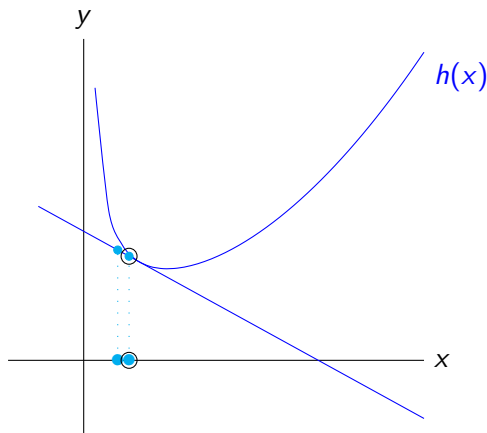
An optimistic algorithm



$$\ell(x_{avg}^\dagger) = \frac{1}{T} \sum_t \ell(x_t^\dagger) \leq \frac{1}{T} \sum_t \ell(x_t^*) = \ell(x^*) \leq h(x^*)$$

Upper bound on regret: $h(x_{avg}^\dagger) - \ell(x_{avg}^\dagger)$

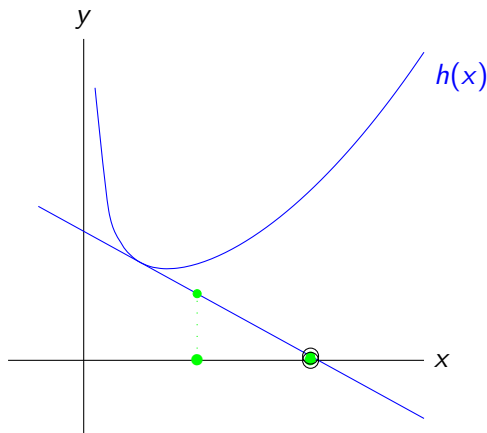
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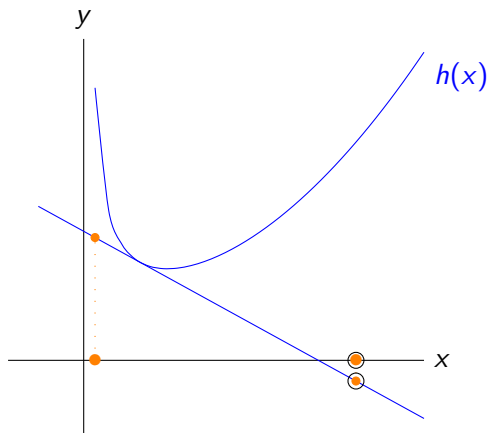
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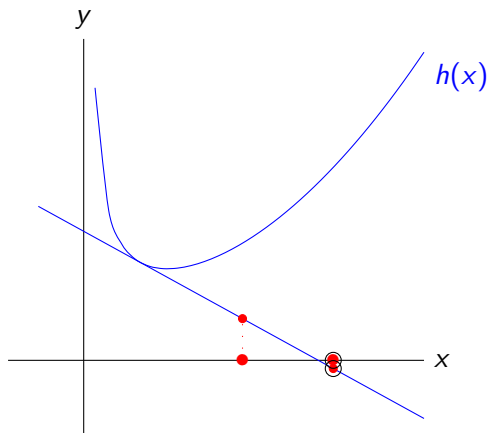
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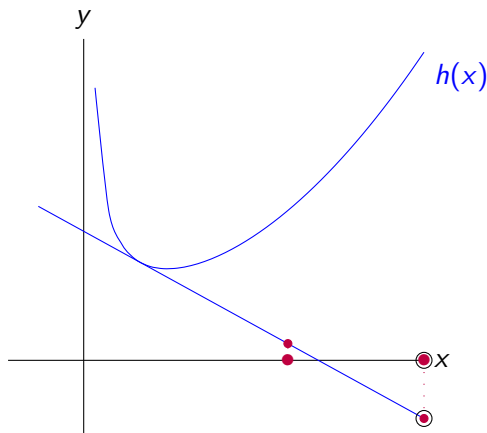
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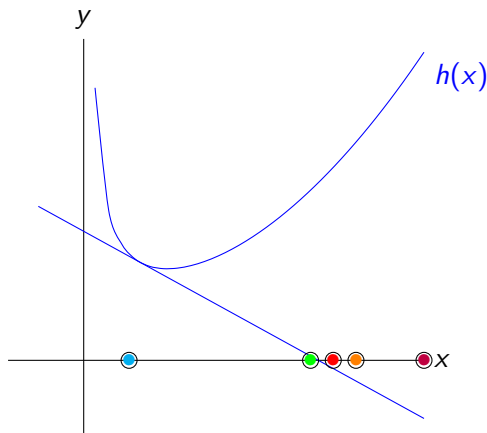
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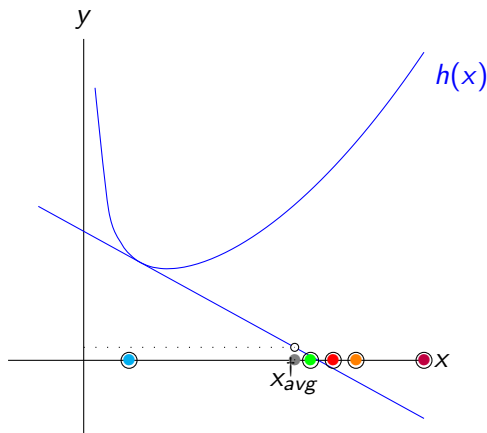
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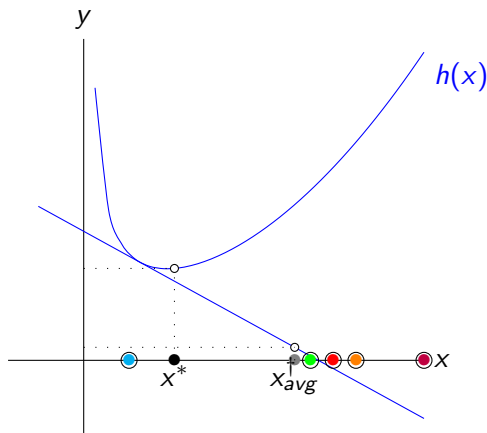
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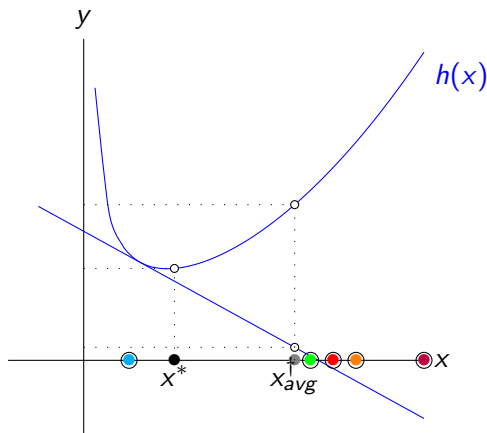
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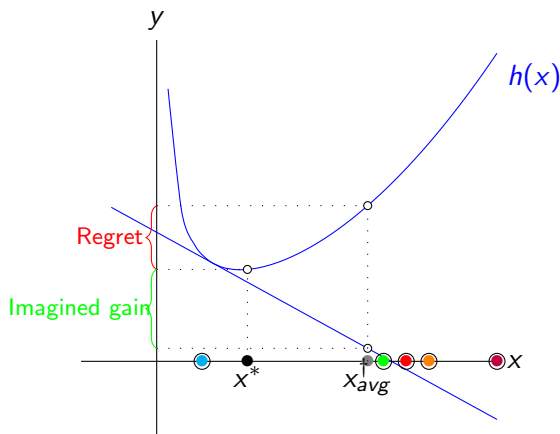
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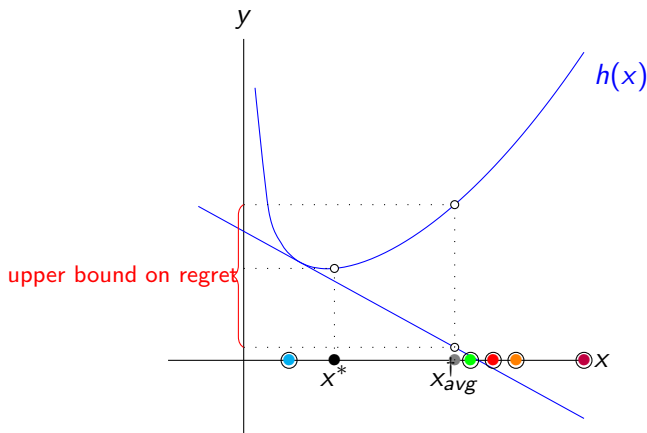
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$$\ell(x_{avg}^\dagger) = \frac{1}{T} \sum_t \ell(x_t^\dagger) \leq \frac{1}{T} \sum_t \ell(x_t^*) = \ell(x^*) \leq h(x^*)$$

Upper bound on regret: $h(x_{avg}^\dagger) - \ell(x_{avg}^\dagger)$

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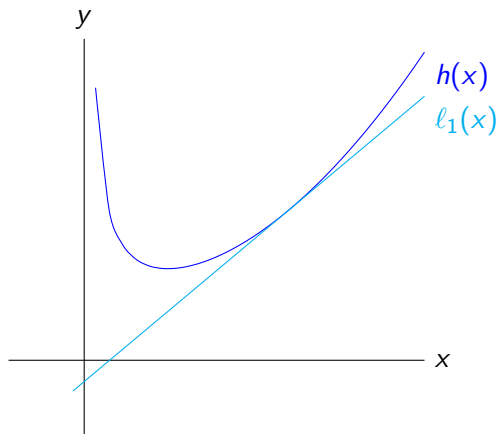


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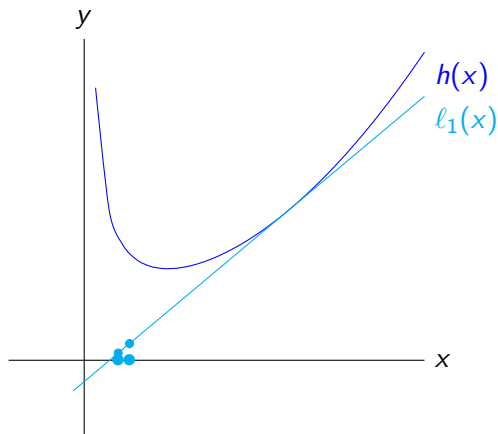
- ▶ If x_{avg}^\dagger was known, tangent at this point would be a linear function with 0 gap: $\ell(x_{avg}^\dagger) = h(x_{avg}^\dagger)$
- ▶ At time t , use current average as a guess for x_{avg}^\dagger and take tangent (slope is gradient) at that point.
- ▶ Algorithm that uses a different tangent at every step.

Algorithm



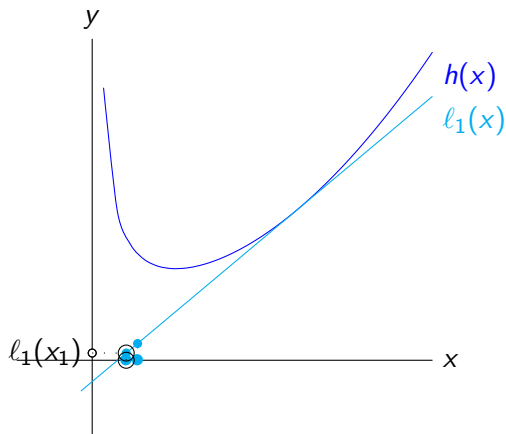
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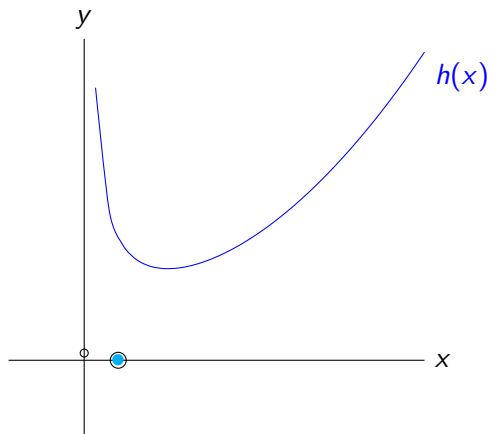
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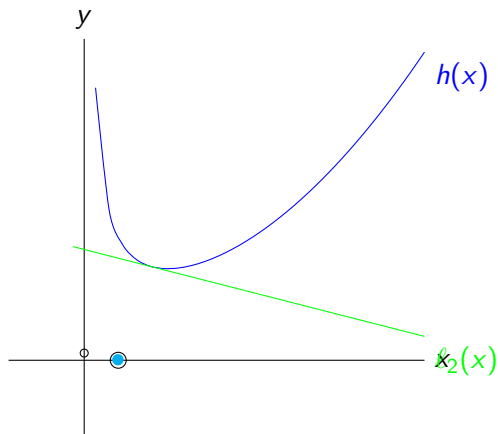
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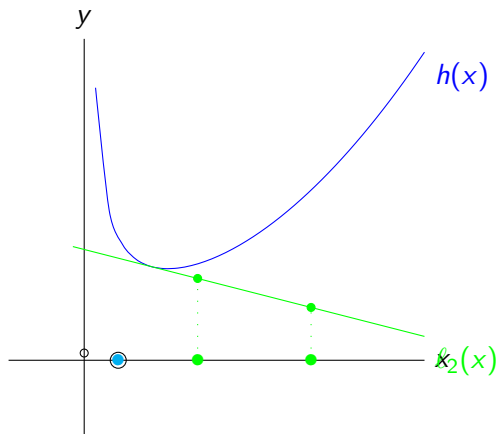
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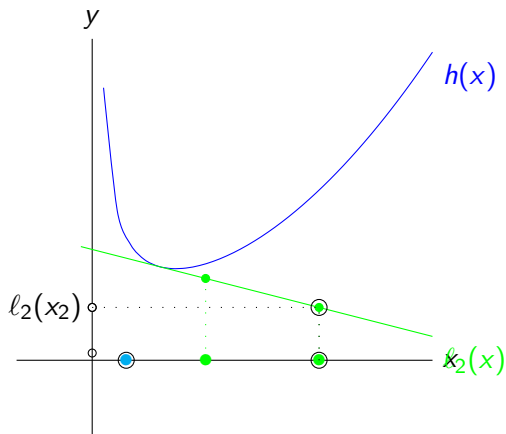
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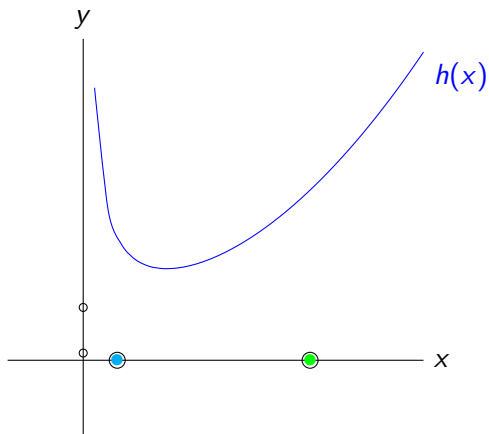
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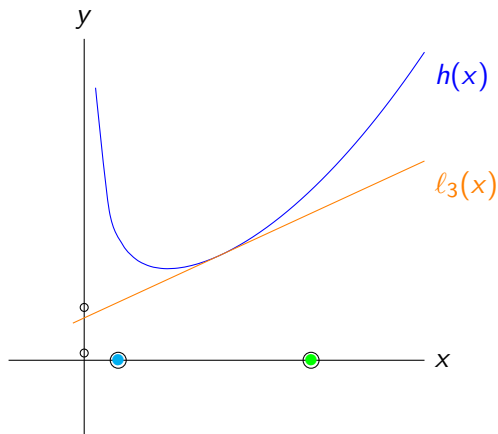
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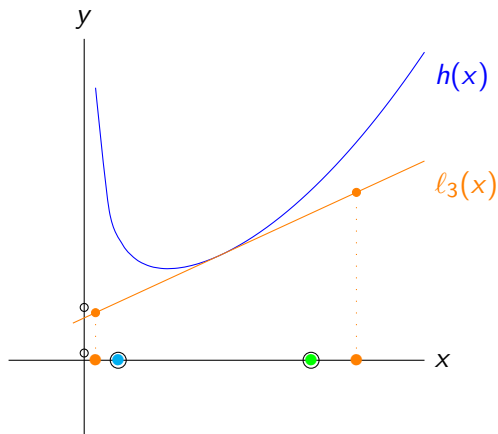
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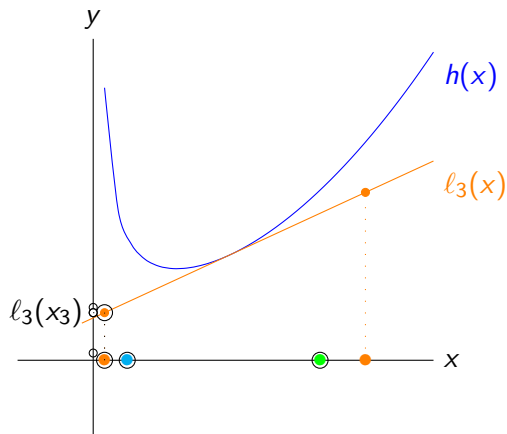
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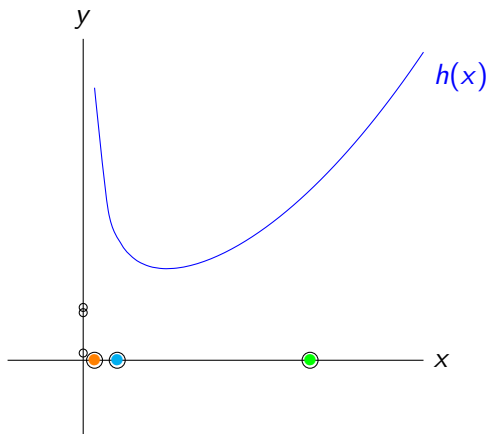
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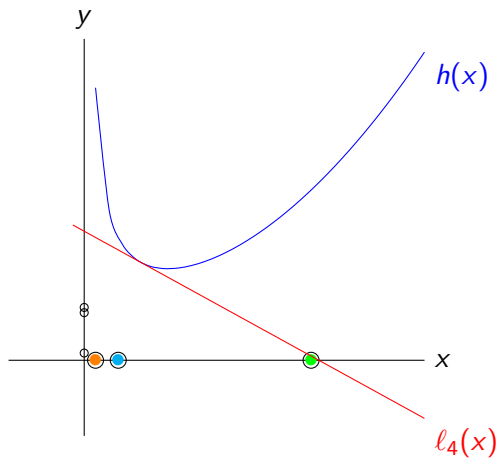
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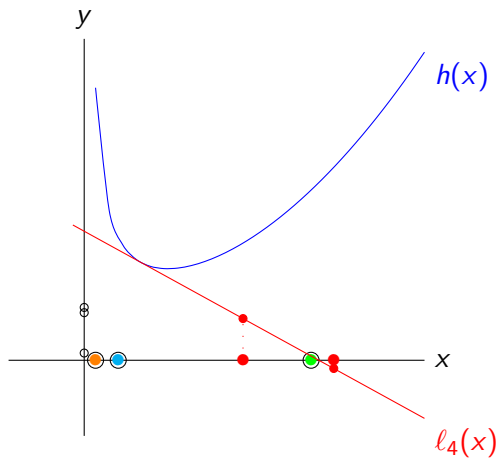
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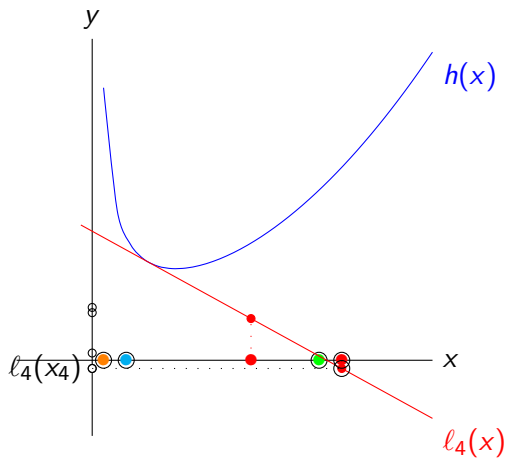
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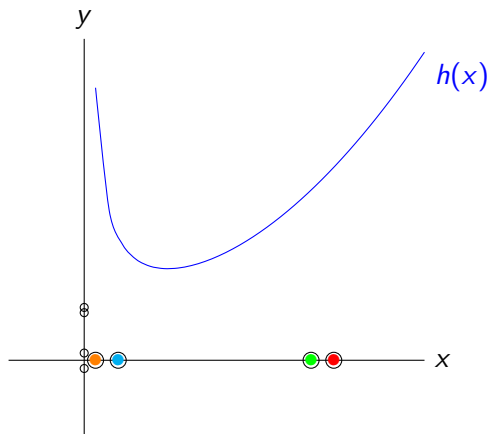
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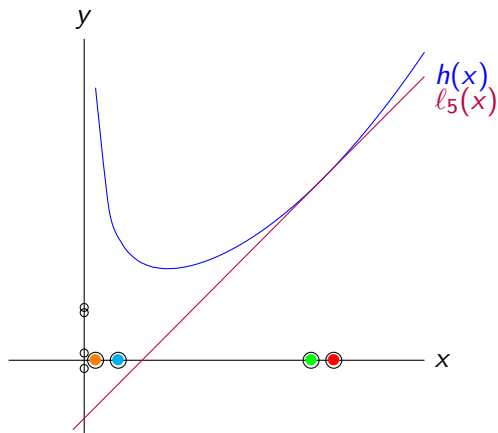
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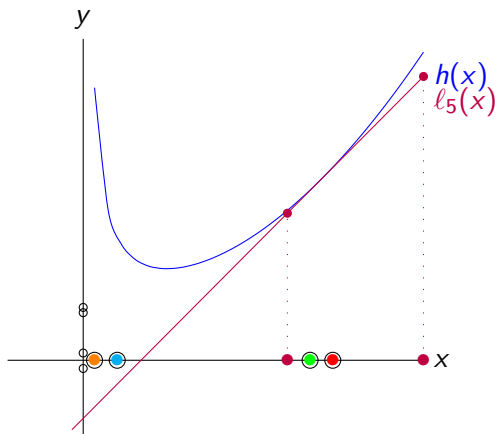
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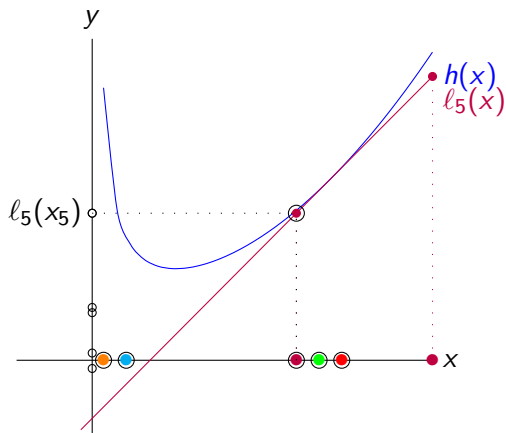
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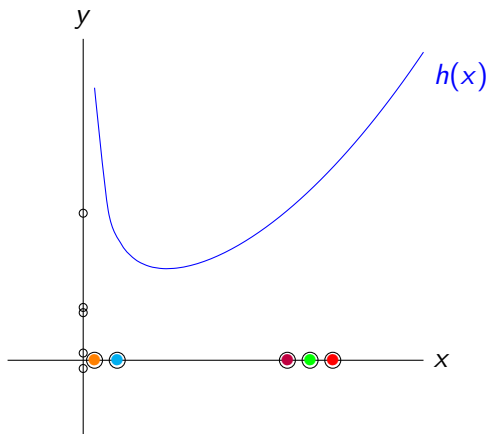
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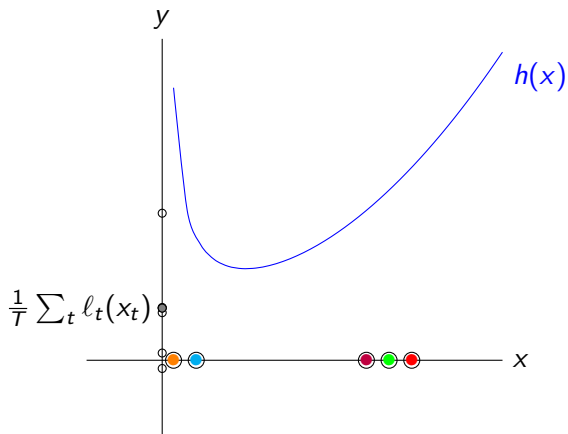
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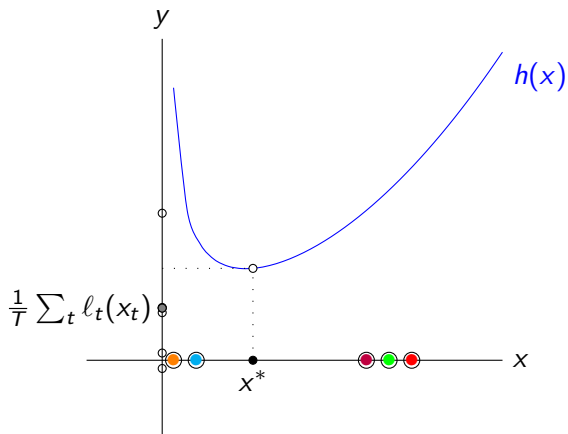
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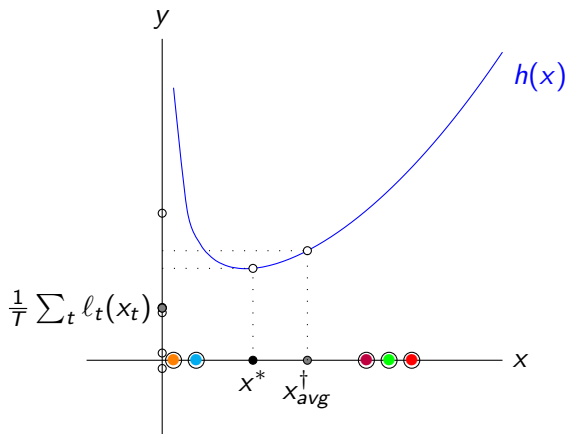
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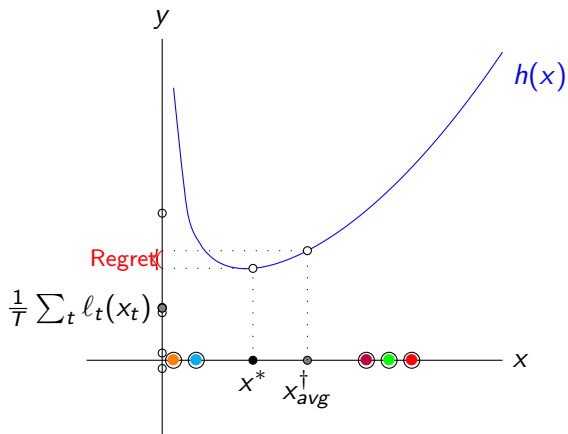
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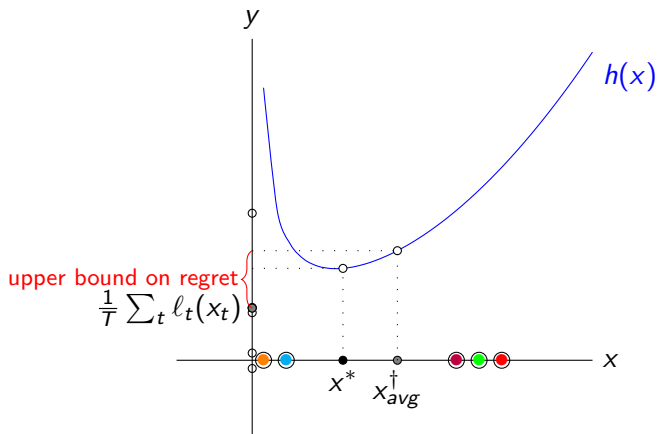
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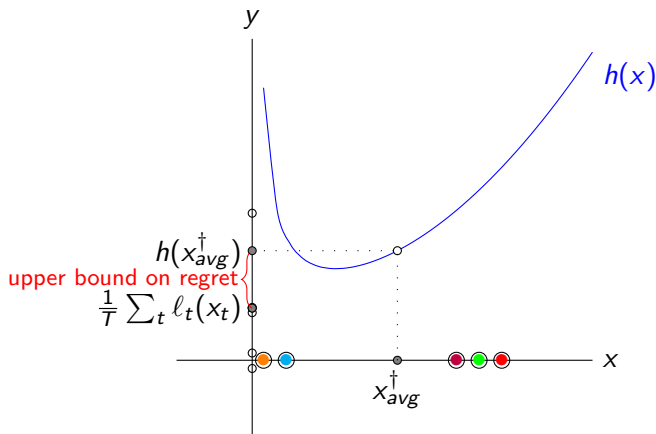
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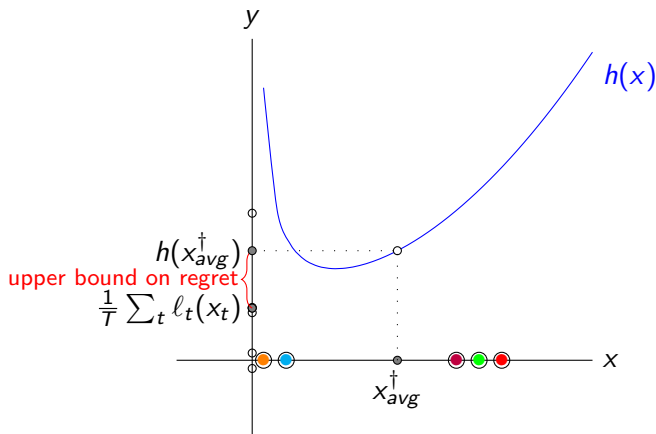
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- ▶ Summing up for $t = 1, \dots, T$ gives $O(\frac{\beta}{T})$ bound on $h(x_{avg}) - \frac{1}{T} \sum_t \ell_t(x_t^\dagger)$. For convex but non-smooth functions, bound degrades to $\tilde{O}(1/\sqrt{T})$.

Algorithm Outline

Algorithm 1 Algorithm for minimizing $h(\frac{1}{T} \sum_{t=1}^T W\mathbf{x}_{t,a_t})$, with known W .

for all $t = 1 \dots T$ **do**

Observe $\mathbf{x}_{t,a}$ for all $a \in A_t$.

Guess $\ell_t(\cdot)$.

$$a_t := \arg \min_{a \in A_t} \ell_t(W\mathbf{x}_{t,a}).$$

end for

Note:

- ▶ Optimistic guess: $\ell_t(W\mathbf{x}_{t,a})$ lower bounds $h(W\mathbf{x}_{t,a})$
- ▶ Regret bounded by the gap at played arms:

$$h\left(\frac{1}{T} \sum_t W\mathbf{x}_{t,a_t}\right) - \frac{1}{T} \sum_t \ell_t(W\mathbf{x}_{t,a_t}) \leq \tilde{O}\left(\frac{\log(d)}{\sqrt{T}}\right)$$

Handling unknown W

Replace W by its optimistic estimate: in this case lower confidence bound.

Algorithm 2 Algorithm for unknown W

for all $t = 1 \dots T$ **do**

Observe $\mathbf{x}_{t,a}$ for all $a \in A_t$.

For all $a \in A_t$, compute lower confidence bound (LCB) $\tilde{W}_{t,a}$ as in linear contextual MAB.

Guess tangent $\ell_t(\cdot)$.

Play arm

$$a_t := \arg \min_{a \in A_t} \ell_t(\tilde{W}_{t,a} \mathbf{x}_{t,a}).$$

Observe $\mathbf{v}_t := \mathbf{v}_{t,a_t}$, with expected value $W \mathbf{x}_{t,a_t}$

end for

Additional term added to regret:

$$\left(\frac{1}{T} \sum_t \ell_t(W \mathbf{x}_{t,a_t}) - \frac{1}{T} \sum_t \ell_t(\tilde{W} \mathbf{x}_{t,a_t}) \right) \leq \tilde{O}\left(\frac{d}{\sqrt{T}}\right)$$

Further difficulties

So far: algorithm for minimizing a convex function on average decision. How to handle “maximize concave function given constraint set S ”

- ▶ “Constraints only” case can be handled by posing problem as “minimize distance from the constraint set”
- ▶ If OPT known, convert objective into constraint.
- ▶ Estimating OPT, requires further exploration, incurring suboptimal regret $dT^{-1/3}$
- ▶ Getting d/\sqrt{T} regret (or a tighter lower bound) is open

Thank You