# Multi-agent Team Formation Diversity Beats Strength? 

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## Appendix

## A Diversity Beats Strength

In Section 3.1 we presented an example with nondeterministic agents that showed that a diverse team can play better than a uniform team made of copies of the strongest agent. The full description of the agents used in the example can be seen in Table 1, where we show the pdf of the agents for each world state. We considered the utility vector $<1,0,0>$ for all world states.

## B Optimal Voting Rules

We present the derivation of the optimal voting rule, stated in Section 3.1. If an agent $\phi_{i}$ votes for a certain action $a_{x}$, the probability of $a_{x}$ being the action with rank $r$ will be given by $p_{r}^{i}$, the probability that $\phi_{i}$ voted for the action with rank $r$. This will only be true if we assume a uniform prior probability for the ranking of all actions ${ }^{1}$.

Given a certain voting pattern, each possible ranking for the actions in the voting pattern is a mutually exclusive event. Therefore, we sum over all possible ranking combinations where $a_{x}$ is the best action. The votes of each agent are independent, hence the multiplication of the probabilities, as presented in the paper.

Now, we present here the full proof of Theorem 2, stated in Section 3.1.

Proof of Theorem 2 By Assumption 2 we know that we are looking for a tie-breaking rule, as the action chosen by the majority of the votes should always be taken. Let's consider the sets and the voting result described in the Assumption 1. Let $<p_{1}, \ldots, p_{k}>$ be the pdf of agent $\phi_{\text {best }, j}^{\prime}$, and the pdf of the other agents of the subset be $<p_{1}-\epsilon_{i}, p_{2}+\gamma_{i 2}, \ldots, p_{k}+$ $\gamma_{i k}>, \gamma_{i l} \geq 0 \forall l \in(2, k)$ and $\sum_{l=2}^{k} \gamma_{i l}=\epsilon_{i}$. Let $b$ be a rank in $(2, k)$. The probability of $a_{x}$ being the best action is given by:

$$
P_{1}=\left(p_{1}\right) \prod_{i \in W e a k}\left(p_{1}-\epsilon_{i}\right) \prod_{t \in \text { Strong }}\left(p_{b}+\gamma_{t b}\right)
$$

While the probability that $a_{y}$ is the best action is given by:

$$
P_{2}=\left(p_{b}\right) \prod_{i \in \text { Weak }}\left(p_{b}+\gamma_{i b}\right) \prod_{t \in \text { Strong }}\left(p_{1}-\epsilon_{t}\right)
$$

[^0]By Assumption 1, we have that $P_{1}>P_{2}$. We can generate another voting pattern by making one agent $\phi_{\text {weak }}$ in Weak vote for $a_{y}$ and one agent $\phi_{\text {strong }}$ in Strong vote for $a_{x}$. The probability of $a_{x}$ being the best action will change to:

$$
P_{1}^{\prime}=P_{1} * \frac{\left(p_{1}-\epsilon_{\text {strong }}\right)\left(p_{b}+\gamma_{\text {weak }, b}\right)}{\left(p_{1}-\epsilon_{\text {weak }}\right)\left(p_{b}+\gamma_{\text {strong }, b}\right)}
$$

While the probability of $a_{y}$ being the best action will change to:

$$
P_{2}^{\prime}=P_{2} * \frac{\left(p_{1}-\epsilon_{\text {weak }}\right)\left(p_{b}+\gamma_{\text {strong }, b}\right)}{\left(p_{1}-\epsilon_{\text {strong }}\right)\left(p_{b}+\gamma_{\text {weak }, b}\right)}
$$

As $\left(p_{1}-\epsilon_{\text {strong }}\right)>\left(p_{1}-\epsilon_{\text {weak }}\right)$ and $\left(p_{b}+\gamma_{\text {weak }, b}\right)>$ $\left(p_{b}+\gamma_{\text {strong }, b}\right)$ by the Assumption 1, we have that $P_{1}^{\prime}>$ $P_{1}$. Similarly, as $\left(p_{1}-\epsilon_{\text {weak }}\right)<\left(p_{1}-\epsilon_{\text {strong }}\right)$ and $\left(p_{b}+\right.$ $\left.\gamma_{\text {strong }, b}\right)<\left(p_{b}+\gamma_{\text {weak }, b}\right)$ by the Assumption 1, we have that $P_{2}^{\prime}<P_{2}$.

Therefore, assuming that $P_{1}>P_{2}$, we have that $P_{1}^{\prime}>$ $P_{2}^{\prime}$. Hence, for all modifications that can be generated by switching one of the agents, it is better to break ties in favor of the strongest agent. We can use all these voting patterns as a base and apply the same process recursively, to generate all possible voting patterns with a tie. Therefore, it will always be better to break ties in favor of the strongest agent.

Now we consider voting patterns with a tie between more than two options. Let's suppose that in this case breaking ties in favor of the strongest agent $\left(\phi_{b e s t, j}^{\prime}\right)$ is not the optimal voting rule. Therefore, we should break the tie in favor of some option $a_{y}$. This implies that $a_{y}$ has a higher probability of being the best action than $a_{x}$, the option chosen by the best agent. Now let's remove the agents that voted in all other options except $a_{x}$ and $a_{y}$. This affects the probability of $a_{x}$ and $a_{y}$ being the best action in the same way. Therefore, we should still break ties in favor of option $a_{y}$. However, we already showed that when there are two options we should break ties in favor of the strongest agent. Hence, we should break the tie in favor of option $a_{x}$. So, by contradiction, we see that if there is a tie between more than two options we should still break ties in favor of the strongest agent.

If the strongest agent of the team is not one of the agents involved in the tie, we can ignore the opinion of the strongest agent according to Assumption 2, and break the tie in favor of the strongest agent from the ones involved in the tie, because Assumption 1 applies to any subset of the agents.

| Agent | State 1 | State 2 | State 3 | State 4 | Strength |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Agent 1 | $<0.99,0.01,0>$ | $<0,0.99,0.01>$ | $<0.99,0,0.01>$ | $<0.99,0.01,0>$ | 0.7425 |
| Agent 2 | $<0,0.99,0.01>$ | $<0.99,0.01,0>$ | $<0.99,0,0.01>$ | $<0,0.01,0.99>$ | 0.4950 |
| Agent 3 | $<0.99,0.005,0.005>$ | $<0.99,0.005,0.005>$ | $<0,0.5,0.5>$ | $<0,0.5,0.5>$ | 0.4950 |
| Agent 4 | $<0.99,0.01,0>$ | $<0.99,0.004,0.006>$ | $<0,0.4,0.6>$ | $<0.99,0.003,0.007>$ | 0.7425 |
| Agent 5 | $<0,0.3,0.7>$ | $<0,0.7,0.3>$ | $<0.99,0.005,0.005>$ | $<0.99,0.002,0.008>$ | 0.4950 |

Table 1: A team of non-deterministic agents that can overcome copies of the best agent.

## C Definition of Fuego $\Delta$ and Fuego $\Theta$

In Section 4.2 we used agents Fuego $\Delta$ and Fuego $\Theta$ in our experiments. We present here the description of these agents. Fuego follows an UCT Monte Carlo Go algorithm, so it uses heuristics to simulate games during the Monte Carlo Simulations. There are mainly 5 possible heuristics in Fuego's code. These heuristics have a hierarchical order, and the original Fuego agent follows the order $<$ Atari Capture, Atari Defend, Lowlib, Pattern > (The heuristic called Nakade is not enabled by default). We created a variation of Fuego, that will be called Fuego $\Delta$, that follows the order $<$ Atari Defend, Atari Capture, Pattern, Nakade, Lowlib>. We also created Fuego $\Theta$, that follows the order $<$ Atari Defend, Nakade, Pattern, Atari Capture, Lowlib>. The memory available for Fuego $\Delta$ and Fuego $\Theta$ is half of the memory available for Fuego.


[^0]:    ${ }^{1}$ The appendix will be updated with a more formal explanation.

